

Regular machine: global environments

%datatype *clos*

%datatype *c-env*

%datatype *k-env*

%datatype *stack*

%name *clos* *c*

%name *c-env* \mathcal{E}

%name *k-env* \mathcal{E}_μ

%name *stack* \mathcal{S}

c ::= $(t, \mathcal{E}, \mathcal{E}_\mu)$

\mathcal{E}_μ ::= ()
| $(\mathcal{S}; \mathcal{E}_\mu)$

\mathcal{E} ::= ()

| $(c; \mathcal{E})$

\mathcal{S} ::= []
| $c :: \mathcal{S}$

%datatype *state*

%name *state* σ

σ ::= $\langle t, \mathcal{E}, \mathcal{E}_\mu, \mathcal{S} \rangle$

Regular machine: evaluation rules

%judgment $\sigma_1 \rightsquigarrow \sigma_2$

$\langle n, \mathcal{E}, \mathcal{E}_\mu, \mathcal{S} \rangle \rightsquigarrow \langle t, \mathcal{E}', \mathcal{E}'_\mu, \mathcal{S} \rangle$	[k·var] when $\mathcal{E}(n) = (t, \mathcal{E}', \mathcal{E}'_\mu)$
$\langle (t u), \mathcal{E}, \mathcal{E}_\mu, \mathcal{S} \rangle \rightsquigarrow \langle t, \mathcal{E}, \mathcal{E}_\mu, (u, \mathcal{E}, \mathcal{E}_\mu) :: \mathcal{S} \rangle$	[k·app]
$\langle \lambda t, \mathcal{E}, \mathcal{E}_\mu, c :: \mathcal{S} \rangle \rightsquigarrow \langle t, (c; \mathcal{E}), \mathcal{E}_\mu, \mathcal{S} \rangle$	[k·abs]
$\langle \text{catch } t, \mathcal{E}, \mathcal{E}_\mu, \mathcal{S} \rangle \rightsquigarrow \langle t, \mathcal{E}, (\mathcal{S}; \mathcal{E}_\mu), \mathcal{S} \rangle$	[k·catch]
$\langle \text{throw } \alpha t, \mathcal{E}, \mathcal{E}_\mu, \mathcal{S} \rangle \rightsquigarrow \langle t, \mathcal{E}, \mathcal{E}_\mu, \mathcal{S}' \rangle$	[k·throw] when $\mathcal{E}_\mu(\alpha) = \mathcal{S}'$

%unique $+ \sigma_1 \rightsquigarrow -1 \sigma_2$