

1 First simulation (soundness)

1.1 Syntax

1.1.1 Index, indices and tables

%datatype *index* $n \ \alpha$
 $n ::= 0$
 $\quad \mid n+1$

%datatype *vector*
%name *vector* Z
 $Z ::= []$
 $\quad \mid n :: Z$

%datatype *table*
%name *table* Z_n
 $Z_n ::= []$
 $\quad \mid Z :: Z_n$

1.1.2 Term

%datatype *term*
%name *term* t
 $t ::= n$
 $\quad \mid t_1 t_2$
 $\quad \mid \lambda t$
 $\quad \mid \text{catch } t$
 $\quad \mid \text{throw } \alpha t$

Remark. Syntax of safe λ_{C} -terms:

$t ::= n$
 $\quad \mid t_1 t_2$
 $\quad \mid \lambda t$
 $\quad \mid \text{get-context } t$
 $\quad \mid \text{set-context } \alpha t$

1.2 Subtraction

%judgment $n_1 \dot{-} n_2 = n_3$

$n_1 \dot{-} 0 = n_1$ [minus0]
 $(n_1 + 1) \dot{-} (n_2 + 1) = n_3$ [minus1] when $n_1 \dot{-} n_2 = n_3$

%mode $+n_1 \dot{-} +n_2 = -n_3$

%worlds $() \ n_1 \dot{-} n_2 = n_3$

%terminates $(n_1) \ n_1 \dot{-} n_2 = n_3$

%unique $+n_1 \dot{-} +n_2 = -n_3$

%lemma $\forall n: \text{index} \cdot n \dot{-} n = 0$ [minus0f]

Proof.

$$\frac{0: \text{index} \cdot 0 \dot{-} 0 = 0 \text{ [minus0]} \quad \text{[minus0f]} [k1]}{n: \text{index} \cdot n \dot{-} n = 0 \text{ [minus0f]}} \text{ [k2]}$$

$$\frac{D}{n+1: \text{index} \cdot (n+1) \dot{-} (n+1) = 0 \text{ [minus0f]}} \text{ [minus0f]}$$

%mode $+n \cdot -D$ [minus0f]

%worlds $() \ n \cdot -D$ [minus0f]

%total $(\alpha) \ n \cdot -D$ [minus0f]

%lemma $n_1 \dot{-} (n_2 + 1) = n_3 \Rightarrow n_1 \dot{-} n_2 = (n_3 + 1)$ [minus0swc]

Proof.

$$\frac{D}{(n_1 + 1) \dot{-} (0 + 1) = n_3 \Rightarrow (n_1 + 1) \dot{-} 0 = (n_3 + 1) \text{ [minus0swc]}} \text{ [k3]}$$

$$\frac{\frac{D_1}{n_1 \dot{-} (n_2 + 1) = n_3} \Rightarrow \frac{D_2}{n_1 \dot{-} n_2 = (n_3 + 1)} \text{ [minus0swc]}}{\frac{D_1}{(n_1 + 1) \dot{-} ((n_2 + 1) + 1) = n_3} \Rightarrow \frac{D_2}{(n_1 + 1) \dot{-} (n_2 + 1) = (n_3 + 1)} \text{ [minus0f]}} \text{ [k2]}$$

%mode $+D_1 \Rightarrow -D_2$ [minus0swc]

%worlds $() \ D_1 \Rightarrow D_2$ [minus0swc]

%total $D_1 \ D_1 \Rightarrow D_2$ [minus0swc]

%lemma $n_1 \dot{-} n_2 = n_3 \Rightarrow n_1 \dot{-} n_2 = n_3$ [minus0swc]

Proof.

$$\frac{n_1: \text{index} \cdot n_1 \dot{-} n_1 = 0 \text{ [minus0f]}}{n_1 \dot{-} 0 = n_1 \text{ [minus0]} \Rightarrow n_1 \dot{-} n_1 = 0 \text{ [minus0swc]}} \text{ [k1]}$$

$$\frac{\frac{D_1}{n_1 \dot{-} n_2 = n_3} \Rightarrow \frac{D_2}{n_1 \dot{-} n_2 = n_2} \text{ [minus0swc]}}{\frac{D_1}{(n_1 + 1) \dot{-} (n_2 + 1) = n_3} \Rightarrow \frac{D_2}{(n_1 + 1) \dot{-} n_2 = (n_2 + 1)} \text{ [minus0swc]}} \text{ [k2]}$$

$$\frac{\frac{D_1}{n_1 \dot{-} n_2 = n_3} \Rightarrow \frac{D_2}{n_1 \dot{-} n_2 = n_2} \text{ [minus0swc]}}{\frac{D_1}{(n_1 + 1) \dot{-} (n_2 + 1) = n_3} \Rightarrow \frac{D_2}{(n_1 + 1) \dot{-} n_2 = (n_2 + 1)} \text{ [minus0swc]}} \text{ [k2]}$$

%mode $+D_1 \Rightarrow -D_2$ [minus0swc]

%worlds $() \ D_1 \Rightarrow D_2$ [minus0swc]

%total $D_1 \ D_1 \Rightarrow D_2$ [minus0swc]

1.2.1 Fetch (indices)

%judgment $Z(n_1) = n_2$

$(n :: Z)(0) = n$ [fetch0f]

$(n :: Z)(n_1 + 1) = n_2$ [fetch1f] when $Z(n_1) = n_2$

%mode $+Z(n_1) = -n_2$

%worlds $() \ Z(n_1) = n_2$

%terminates $n_1 \ Z(n_1) = n_2$

%unique $+Z(n_1) = -n_2$

1.2.2 Fetch (table)

%judgment $Z_n(\alpha) = Z$

$(Z :: Z_n)(0) = Z$ [fetch0f]

$(Z :: Z_n)(\alpha + 1) = Z'$ [fetch1f] when $Z_n(\alpha) = Z$

%mode $+Z_n(\alpha) = -Z$

%worlds $() \ Z_n(\alpha) = Z$

%terminates $\alpha \ Z_n(\alpha) = Z$

%unique $+Z_n(\alpha) = -Z$

1.2.3 Compute

%judgment $n_1 \dot{-} Z(n_2) = n_3$

$n \dot{-} Z(t) = g$ [compute] when $Z(t) = k$, $n \dot{-} k = g$

%mode $+n_1 \dot{-} +Z(n_2) = -n_3$

%worlds $() \ n_1 \dot{-} Z(n_2) = n_3$

%terminates $() \ n_1 \dot{-} Z(n_2) = n_3$

%unique $+n_1 \dot{-} +Z(n_2) = -n_3$

2 Safe λ_{C} -terms

2.1 Safety

%judgment $n \in Z$

$n \in (n :: Z)$ [members]

$n \in (n' :: Z)$ [members] when $n \in Z$

%mode $+n \in +Z$

%worlds $() \ n \in Z$

%terminates $Z \ n \in Z$

%lemma $Z(n) = k \Rightarrow k \in Z$ [target]

Proof.

$$\frac{(n :: Z)(0) = n \text{ [fetch0f]} \Rightarrow n \in (n :: Z) \text{ [members]} \text{ [target]} [k1]}{\frac{D_1}{Z(n) = k} \Rightarrow \frac{D_2}{k \in Z} \text{ [target]}} \text{ [k2]}$$

$$\frac{\frac{D_1}{Z(n) = k} \Rightarrow \frac{D_2}{k \in Z} \text{ [target]}}{\frac{D_1}{(Z' :: Z)(\alpha + 1) = k} \Rightarrow \frac{D_2}{k \in (Z' :: Z)} \text{ [members]} \text{ [target]}} \text{ [k2]}$$

%mode $+D_1 \Rightarrow -D_2$ [target]

%worlds $() \ D_1 \Rightarrow D_2$ [target]

%total $(D_1) \ D_1 \Rightarrow D_2$ [target]

%judgment $\text{Safe}_n^{Z, Z'}(t)$

$\text{Safe}_n^{Z, Z'}(g)$ [safe] when $n \dot{-} g = k$, $k \in Z$

$\text{Safe}_n^{Z, Z'}(u)$ [safe] when $\text{Safe}_n^{Z, Z'}(t)$, $\text{Safe}_n^{Z, Z'}(u) = u'$

$\text{Safe}_n^{Z, Z'}(\lambda t)$ [safe] when $\text{Safe}_n^{Z, Z'}(t)$

$\text{Safe}_n^{Z, Z'}(\text{catch } t)$ [safe] when $\text{Safe}_n^{Z, Z'}(t)$

$\text{Safe}_n^{Z, Z'}(\text{throw } \alpha t)$ [safe] when $Z_n(\alpha) = Z'$, $\text{Safe}_n^{Z, Z'}(t)$

%mode $\text{Safe}_n^{Z, Z'}(+t) = -t'$

%worlds $() \ \text{Safe}_n^{Z, Z'}(t) = t'$

%terminates $t \ \text{Safe}_n^{Z, Z'}(t)$

2.2 From local indices to global indices

%judgment $\dot{-}^Z, Z'(t_1) = t_2$

$\dot{-}^Z, Z'(t) = g$ [l1] when $n \dot{-} Z(t) = g$

$\dot{-}^Z, Z'(t u) = t' u'$ [l2] when $\dot{-}^Z, Z'(t) = t'$, $\dot{-}^Z, Z'(u) = u'$

$\dot{-}^Z, Z'(\lambda t) = \lambda t'$ [l3] when $\dot{-}^Z, Z'(t) = t'$

$\dot{-}^Z, Z'(\text{get-context } t) = \text{catch } t'$ [l4] when $\dot{-}^Z, Z'(t) = t'$

$\dot{-}^Z, Z'(\text{set-context } \alpha t) = \text{throw } \alpha t'$ [l5] when $Z_n(\alpha) = Z'$, $\dot{-}^Z, Z'(t) = t'$

%mode $\dot{-}^Z, Z'(+t) = -t'$

%worlds $() \ \dot{-}^Z, Z'(t) = t'$

%terminates $t \ \dot{-}^Z, Z'(t) = t'$

%unique $\dot{-}^Z, Z'(+t) = -t'$

Lemma $\downarrow^Z X_n(t) = t' \Rightarrow \text{Safe}_{\alpha}^Z X_n(t')$ [1+safe]

Proof.

$$\frac{\frac{\frac{D_{11}}{Z(t) = k} \Rightarrow \frac{D_2}{k \in Z} \text{ [target]}}{\frac{D_{12}}{n \perp k = g} \Rightarrow \frac{D_3}{n \perp g = k} \text{ [infix swap]}}}{\frac{D_{11}}{Z(t) = k} \frac{D_{12}}{n \perp k = g} \text{ [compose]}}{\frac{\downarrow^Z X_n(t) = g}{\downarrow^Z X_n(t) = g} \text{ [1]}} \Rightarrow \frac{\frac{D_2}{n \perp g = k} \frac{D_3}{k \in Z}}{\text{Safe}_{\alpha}^Z X_n(g)} \text{ [sub]} \text{ [1+safe]} \quad [k1]$$

$$\frac{\frac{\frac{D_{11}}{\downarrow^Z X_n(t) = t'} \Rightarrow \frac{D_1}{\text{Safe}_{\alpha}^Z X_n(t')} \text{ [1+safe]}}{\frac{D_{12}}{\downarrow^Z X_n(u) = u'} \Rightarrow \frac{D_2}{\text{Safe}_{\alpha}^Z X_n(u')} \text{ [1+safe]}}}{\frac{D_{11}}{\downarrow^Z X_n(t) = t'} \frac{D_{12}}{\downarrow^Z X_n(u) = u'} \text{ [1]}} \Rightarrow \frac{\frac{D_1}{\text{Safe}_{\alpha}^Z X_n(t')} \frac{D_2}{\text{Safe}_{\alpha}^Z X_n(u')}}{\text{Safe}_{\alpha}^Z X_n(t'u')} \text{ [sub]} \text{ [1+safe]} \quad [k2]$$

$$\frac{\frac{D_{11}}{\downarrow^{(n+1)Z} X_n(t) = t'} \Rightarrow \frac{D_1}{\text{Safe}_{\alpha}^{(n+1)Z} X_n(t')} \text{ [1+safe]}}{\frac{D_{12}}{\downarrow^{(n+1)Z} X_n(u) = u'} \Rightarrow \frac{D_2}{\text{Safe}_{\alpha}^{(n+1)Z} X_n(u')} \text{ [1+safe]}}}{\frac{D_{11}}{\downarrow^{(n+1)Z} X_n(t) = t'} \frac{D_{12}}{\downarrow^{(n+1)Z} X_n(u) = u'} \text{ [1]}} \Rightarrow \frac{\frac{D_1}{\text{Safe}_{\alpha}^{(n+1)Z} X_n(t')} \frac{D_2}{\text{Safe}_{\alpha}^{(n+1)Z} X_n(u')}}{\text{Safe}_{\alpha}^{(n+1)Z} X_n(t'u')} \text{ [sub]} \text{ [1+safe]} \quad [k3]$$

$$\frac{\frac{D_{11}}{\downarrow^Z X_n(\text{get-context } t) = \text{catch } t'} \Rightarrow \frac{D_1}{\text{Safe}_{\alpha}^Z X_n(\text{get-context } t')} \text{ [1+safe]}}{\frac{D_{12}}{\downarrow^Z X_n(\text{set-context } \alpha t) = \text{throw } \alpha t'} \Rightarrow \frac{D_2}{\text{Safe}_{\alpha}^Z X_n(\text{throw } \alpha t')} \text{ [1+safe]}}}{\frac{D_{11}}{\downarrow^Z X_n(\text{get-context } t) = \text{catch } t'} \frac{D_{12}}{\downarrow^Z X_n(\text{set-context } \alpha t) = \text{throw } \alpha t'} \text{ [1]}} \Rightarrow \frac{\frac{D_1}{\text{Safe}_{\alpha}^Z X_n(\text{get-context } t')} \frac{D_2}{\text{Safe}_{\alpha}^Z X_n(\text{throw } \alpha t')}}{\text{Safe}_{\alpha}^Z X_n(\text{throw } \alpha t')} \text{ [sub]} \text{ [1+safe]} \quad [k4]$$

$$\frac{\frac{D_{11}}{\downarrow^Z X_n(\text{get-context } t) = \text{catch } t'} \Rightarrow \frac{D_1}{\text{Safe}_{\alpha}^Z X_n(\text{get-context } t')} \text{ [1+safe]}}{\frac{D_{12}}{\downarrow^Z X_n(\text{set-context } \alpha t) = \text{throw } \alpha t'} \Rightarrow \frac{D_2}{\text{Safe}_{\alpha}^Z X_n(\text{throw } \alpha t')} \text{ [1+safe]}}}{\frac{D_{11}}{\downarrow^Z X_n(\text{get-context } t) = \text{catch } t'} \frac{D_{12}}{\downarrow^Z X_n(\text{set-context } \alpha t) = \text{throw } \alpha t'} \text{ [1]}} \Rightarrow \frac{\frac{D_1}{\text{Safe}_{\alpha}^Z X_n(\text{get-context } t')} \frac{D_2}{\text{Safe}_{\alpha}^Z X_n(\text{throw } \alpha t')}}{\text{Safe}_{\alpha}^Z X_n(\text{throw } \alpha t')} \text{ [sub]} \text{ [1+safe]} \quad [k5]$$

mode $+D_1 \Rightarrow -D_2$ [1+safe]
worlds $() D_1 \Rightarrow D_2$ [1+safe]
total $(D_1) D_1 \Rightarrow D_2$ [1+safe]

□

2.3 Examples

1 = 0 + 1.
 2 = 1 + 1.
 3 = 2 + 1.

solve $- : \text{Safe}_{\alpha}^{\downarrow}(\lambda(\text{catch } \lambda(0(\text{throw } 0))))$

Remark. This example fails as expected:

solve $- : \text{Safe}_{\alpha}^{\downarrow}(\lambda(\lambda(\text{catch } \lambda(1(\text{throw } 0))))$

solve $- : \downarrow^{\downarrow}(\lambda(\lambda(\text{get-context } \lambda(1(\text{set-context } 0)))) = \lambda(\text{catch } \lambda(1(\text{throw } 0)))$

solve $d_1 : \downarrow^{\downarrow}(\lambda(\lambda(\text{get-context } \lambda(1(\text{set-context } 0)))) = \lambda(\text{catch } \lambda(1(\text{throw } 0)))$

solve $- : d_1 \Rightarrow D_2$ [1+safe]

2.4 Closure, environment and stack

datatype *class* **%name** *class* *c*
datatype *c-env* **%name** *c-env* \mathcal{E}
datatype *k-env* **%name *k-env* \mathcal{E}_p
datatype *stack* **%name** *stack* \mathcal{S}**

$c ::= (t, \mathcal{E}, \mathcal{E}_p)$

$\mathcal{E} ::= ()$
 $| (c; \mathcal{E})$

$\mathcal{E}_p ::= ()$
 $| (\mathcal{S}; \mathcal{E}_p)$

$\mathcal{S} ::= []$
 $| c; \mathcal{S}$

datatype *state*

%name *state* σ

$\sigma ::= (t, \mathcal{E}, \mathcal{E}_p, \mathcal{S})$

2.5 Judgments

2.5.1 Fetch a closure

judgment $\mathcal{E}(n) = c$

$(c; \mathcal{E})(0) = c$ [fetch]
 $(c; \mathcal{E})(n+1) = c$ [fetch]

mode $+\mathcal{E}(+n) = -c$

worlds $() \mathcal{E}(n) = c$

terminates $\mathcal{E} \mathcal{E}(n) = c$

unique $+\mathcal{E}(+n) = -1c$

2.5.2 Fetch a stack

judgment $\mathcal{E}_p(n) = \mathcal{S}$

$(\mathcal{S}; \mathcal{E}_p)(0) = \mathcal{S}$ [fetch]
 $(\mathcal{S}; \mathcal{E}_p)(n+1) = \mathcal{S}$ [fetch]

mode $+\mathcal{E}_p(+n) = -\mathcal{S}$

worlds $() \mathcal{E}_p(n) = \mathcal{S}$

terminates $\mathcal{E}_p \mathcal{E}_p(n) = \mathcal{S}$

unique $+\mathcal{E}_p(+n) = -1\mathcal{S}$

2.5.3 Evaluation rules

judgment $\sigma_1 \mapsto \sigma_2$

$(k, \mathcal{E}, \mathcal{E}_p, \mathcal{S}) \mapsto (t, \mathcal{E}', \mathcal{E}'_p, \mathcal{S})$ [eval] when $\mathcal{E}(k) = (t, \mathcal{E}', \mathcal{E}'_p)$

$(\lambda t. \mathcal{E}_p.c; \mathcal{E}, \mathcal{S}) \mapsto (t, \mathcal{E}, \mathcal{E}_p, \mathcal{S})$ [eval]
 $(\lambda t. \mathcal{E}_p.c; \mathcal{E}, \mathcal{S}) \mapsto (t, \mathcal{E}, \mathcal{E}_p, \mathcal{S})$ [eval]

$(\text{catch } t. \mathcal{E}, \mathcal{E}_p, \mathcal{S}) \mapsto (t, \mathcal{E}, (\mathcal{S}; \mathcal{E}_p), \mathcal{S})$ [catch]
 $(\text{throw } \alpha t. \mathcal{E}, \mathcal{E}_p, \mathcal{S}) \mapsto (t, \mathcal{E}, \mathcal{E}_p, \mathcal{S})$ [throw]

mode $+\sigma_1 \mapsto -\sigma_2$

worlds $() \sigma_1 \mapsto \sigma_2$

unique $+\sigma_1 \mapsto -1\sigma_2$

2.6 Abstract machine for safe λ_{α} -terms (with indirection tables)

2.6.1 Syntax

datatype *class*

datatype *c-env*

datatype *k-env*

datatype *stack*

%name *class* \mathcal{C}

%name *c-env* \mathcal{E}

%name *k-env* \mathcal{E}_p

%name *stack* \mathcal{S}

$\mathcal{C} ::= (t, n, \mathcal{I}, \mathcal{I}_p, \mathcal{E}, \mathcal{E}_p)$

$\mathcal{E} ::= ()$
 $| (c; \mathcal{E})$

$\mathcal{E}_p ::= ()$
 $| (\mathcal{S}; \mathcal{E}_p)$

$\mathcal{S} ::= []$
 $| c; \mathcal{S}$

datatype *state*

%name *state* σ

$\sigma ::= (t, n, \mathcal{I}, \mathcal{I}_p, \mathcal{E}, \mathcal{E}_p, \mathcal{S})$

2.6.2 Fetch a closure

judgment $\mathcal{E}(n) = \mathcal{C}$

$(c; \mathcal{E})(0) = \mathcal{C}$ [fetch]
 $(c; \mathcal{E})(n+1) = \mathcal{C}$ [fetch]

mode $+\mathcal{E}(+n) = -\mathcal{C}$

worlds $() \mathcal{E}(n) = \mathcal{C}$

terminates $\mathcal{E} \mathcal{E}(n) = \mathcal{C}$

unique $+\mathcal{E}(+n) = -1\mathcal{C}$

2.6.3 Fetch a stack

judgment $\mathcal{E}_p(n) = \mathcal{S}$

$(\mathcal{S}; \mathcal{E}_p)(0) = \mathcal{S}$ [fetch]
 $(\mathcal{S}; \mathcal{E}_p)(n+1) = \mathcal{S}$ [fetch]

mode $+\mathcal{E}_p(+n) = -\mathcal{S}$

worlds $() \mathcal{E}_p(n) = \mathcal{S}$

terminates $\mathcal{E}_p \mathcal{E}_p(n) = \mathcal{S}$

unique $+\mathcal{E}_p(+n) = -1\mathcal{S}$

3 Evaluation rules

judgment $\sigma_1 \mapsto \sigma_2$

$(t, n, \mathcal{I}, \mathcal{I}_p, \mathcal{E}, \mathcal{E}_p, \mathcal{S}) \mapsto (t, n', \mathcal{I}', \mathcal{I}'_p, \mathcal{E}', \mathcal{E}'_p, \mathcal{S})$ [eval]
 when $n \perp \mathcal{I}(t) = g$, $\mathcal{E}(g) = (t, n', \mathcal{I}', \mathcal{I}'_p, \mathcal{E}', \mathcal{E}'_p)$

$(t, n, \mathcal{I}, \mathcal{I}_p, \mathcal{E}, \mathcal{E}_p, \mathcal{S}) \mapsto (t, n, \mathcal{I}, \mathcal{I}_p, \mathcal{E}, \mathcal{E}_p, (n, n, \mathcal{I}, \mathcal{I}_p, \mathcal{E}, \mathcal{E}_p); \mathcal{S})$ [eval]

$(\lambda t. n, \mathcal{I}, \mathcal{I}_p, \mathcal{E}, \mathcal{E}_p, \mathcal{S}) \mapsto (t, n+1, (n+1) \perp \mathcal{I}, \mathcal{I}_p, (t; \mathcal{E}), \mathcal{E}_p, \mathcal{S})$ [eval]

$(\text{get-context } t, n, \mathcal{I}, \mathcal{I}_p, \mathcal{E}, \mathcal{E}_p, \mathcal{S}) \mapsto (t, n, \mathcal{I}, (\mathcal{I}; \mathcal{I}_p), \mathcal{E}, (\mathcal{S}; \mathcal{E}_p), \mathcal{S})$ [catch]

$(\text{set-context } \alpha t, n, \mathcal{I}, \mathcal{I}_p, \mathcal{E}, \mathcal{E}_p, \mathcal{S}) \mapsto (t, n, \mathcal{I}, \mathcal{I}_p, \mathcal{E}, \mathcal{E}_p, \mathcal{S})$ [throw]
 when $\mathcal{I}_p(\alpha) = \mathcal{I}'$, $\mathcal{E}_p(\alpha) = \mathcal{S}'$

mode $+\sigma_1 \mapsto -\sigma_2$

worlds $() \sigma_1 \mapsto \sigma_2$

unique $+\sigma_1 \mapsto -1\sigma_2$

4 Translation

judgment $\mathcal{C}^* = c$

judgment $\mathcal{E}^* = \mathcal{E}$

judgment $\mathcal{E}_p^* = \mathcal{E}_p$

judgment $\mathcal{S}^* = \mathcal{S}$

$(t, n, \mathcal{I}, \mathcal{I}_p, \mathcal{E}, \mathcal{E}_p)^* = (t, \mathcal{E}, \mathcal{E}_p)$ [eval] when $\downarrow^Z X_n(t) = u$, $\mathcal{E}^* = \mathcal{E}$, $\mathcal{E}_p^* = \mathcal{E}_p$

